Tree Recursion

Discussion 4: September 24, 2025

If there are fewer than 3 people in your group, merge your group with another group in the room. If your group has 6 or more students, you're welcome to split into two sub-groups and then sync up at the end. If you want two separate Pensieve documents for the two sub-groups, just have one sub-group add 1000 to their group number.

Now switch to Pensieve:

• Everyone: Go to pensieve.co, log in with your @berkeley.edu email, and enter your group number as the room number (which was in the email that assigned you to this discussion). As long as you all enter the same number (any number), you'll all be using a shared document.

Once you're on Pensieve, you don't need to return to this page; Pensieve has all the same content (but more features). If for some reason Penseive doesn't work, return to this page and continue with the discussion.

Attendance

Your TA will come around during discussion to check you in. You can start on the worksheet before being checked in; you don't need to wait for your TA to get started.

If you didn't attend for a good reason (such as being sick), fill out this form (within 2 weeks of your discussion): attendance form

Getting Started

Say your name and your favorite tree (a particular tree or a kind of tree) in honor of today's topic: tree recursion.

Definition: Tree recursive functions are functions that call themselves more than once.

In this discussion, don't use a Python interpreter to run code until you are confident your solution is correct. Figure things out and check your work by *thinking* about what your code will do. Not sure? Talk to your group!

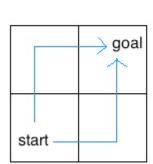
[New] Recursion takes practice. Please don't get discouraged if you're struggling to write recursive functions. Instead, every time you do solve one (even with help or in a group), make note of what you had to realize to make progress. Students improve through practice and reflection.

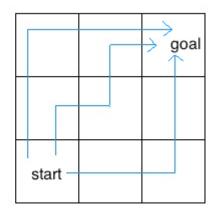
Tree Recursion

For the following questions, don't start trying to write code right away. Instead, start by describing the recursive case in words. Some examples: - In fib from lecture, the recursive case is to add together the previous two Fibonacci numbers. - In double_eights from lab, the recursive case is to check for double eights in the rest of the number. - In count_partitions from lecture, the recursive case is to partition n-m using parts up to size m and to partition n using parts up to size m-1.

Q1: Insect Combinatorics

An insect is inside an m by n grid. The insect starts at the bottom-left corner (1, 1) and wants to end up at the top-right corner (m, n). The insect can only move up or to the right. Write a function paths that takes the height and width of a grid and returns the number of paths the insect can take from the start to the end. (There is a closed-form solution to this problem, but try to answer it with recursion.)





Insect grids.

In the 2 by 2 grid, the insect has two paths from the start to the end. In the 3 by 3 grid, the insect has six paths (only three are shown above).

Hint: What happens if the insect hits the upper or rightmost edge of the grid?

```
def paths(m, n):
   """Return the number of paths from one corner of an
   M by N grid to the opposite corner.
   >>> paths(2, 2)
   >>> paths(5, 7)
   210
   >>> paths(117, 1)
   >>> paths(1, 157)
    0.00
   if m == 1 or n == 1:
        return 1
   return paths(m - 1, n) + paths(m, n - 1)
   # Base case: Look at the two visual examples given. Since the insect
   # can only move to the right or up, once it hits either the rightmost edge
   # or the upper edge, it has a single remaining path -- the insect has
   # no choice but to go straight up or straight right (respectively) at that point.
   # There is no way for it to backtrack by going left or down.
```

The recursive case is that there are paths from the square to the right through an (m, n-1) grid and paths from the square above through an (m-1, n) grid.

Presentation Time: Once your group has converged on a solution, it's time to practice your ability to describe why your recursive case is correct. Nominate someone and have them present to the group for practice. If you want

 ${\it feedback, ask.}$

Tree Recursion with Lists

Some of you already know list operations that we haven't covered yet, such as append. Don't use those today. All you need are list literals (e.g., [1, 2, 3]), item selection (e.g., s[0]), list addition (e.g., [1] + [2, 3]), len (e.g., len(s)), and slicing (e.g., s[1:]). Use those! There will be plenty of time for other list operations when we introduce them next week.

Important: The most important thing to remember about lists is that a non-empty list s can be split into its first element s[0] and the rest of the list s[1:].

```
>>> s = [2, 3, 6, 4]
>>> s[0]
2
>>> s[1:]
[3, 6, 4]
```

Q2: Max Product

Implement max_product, which takes a list of numbers and returns the maximum product that can be formed by multiplying together non-consecutive elements of the list. Assume that all numbers in the input list are greater than or equal to 1.

```
def max_product(s):
    """Return the maximum product of non-consecutive elements of s.
   >>> max_product([10, 3, 1, 9, 2])
                                        # 10 * 9
   90
   >>> max_product([5, 10, 5, 10, 5]) # 5 * 5 * 5
   125
   >>> max_product([])
                                        # The product of no numbers is 1
   1
   if s == []:
        return 1
   if len(s) == 1:
        return s[0]
   else:
        return max(s[0] * max_product(s[2:]), max_product(s[1:]))
        # OR
        return max(s[0] * max_product(s[2:]), s[1] * max_product(s[3:]))
```

First try multiplying the first element by the max_product of everything after the first two elements (skipping the second element because it is consecutive with the first), then try skipping the first element and finding the max_product of the rest. To find which of these options is better, use max.

A great way to get help is to talk to the course staff!

This solution begins with the idea that we either include s[0] in the product or not:

• If we include s[0], we cannot include s[1].

• If we don't include s[0], we can include s[1].

The recursive case is that we choose the larger of: - multiplying s[0] by the max product of s[2:] (skipping s[1]) OR - just the max_product of s[1:] (skipping s[0])

Here are some key ideas in translating this into code: - The built-in max function can find the larger of two numbers, which in this case come from two recursive calls. - In every case, max_product is called on a list of numbers and its return value is treated as a number.

An expression for this recursive case is:

```
max(s[0] * max_product(s[2:]), max_product(s[1:]))
```

Since this expression never refers to s[1], and s[2:] evaluates to the empty list even for a one-element list s, the second base case (len(s) == 1) can be omitted if this recursive case is used.

The recursive solution above explores some options that we know in advance will not be the maximum, such as skipping both s[0] and s[1]. Alternatively, the recursive case could be that we choose the larger of: - multiplying s[0] by the max_product of s[2:] (skipping s[1]) OR - multiplying s[1] by the max_product of s[3:] (skipping s[0] and s[2])

An expression for this recursive case is:

```
\max(s[0] * \max_{product(s[2:])}, s[1] * \max_{product(s[3:])})
```

Description Time: Now try to complete this sentence together: "The recursive case is to choose the larger of _____ and ." When you're done, see how your answer compares to ours.

The recursive case is to choose the larger of the largest product that includes the first but not the second element and the largest product that does not include the first element.

Q3: Sum Fun

Implement sums(n, m), which takes a total n and maximum m. It returns a list of all lists: 1. that sum to n, 2. that contain only positive numbers up to m, and 3. in which no two adjacent numbers are the same.

Two lists with the same numbers in a different order should both be returned.

Here's a recursive approach that matches the template below: build up the result list by building all lists that sum to n and start with k, for each k from 1 to m. For example, the result of sums(5, 3) is made up of three lists: -[[1, 3, 1]] starts with 1, - [[2, 1, 2], [2, 3]] start with 2, and - [[3, 2]] starts with 3.

Hint: Use [k] + s for a number k and list s to build a list that starts with k and then has all the elements of s.

```
>>> k = 2
>>> s = [4, 3, 1]
>>> [k] + s
[2, 4, 3, 1]
```

```
def sums(n, m):
   """Return lists that sum to n containing positive numbers up to m that
   have no adjacent repeats.
   >>> sums(5, 1)
    >>> sums(5, 2)
    [[2, 1, 2]]
   >>> sums(5, 3)
    [[1, 3, 1], [2, 1, 2], [2, 3], [3, 2]]
   >>> sums(5, 5)
    [[1, 3, 1], [1, 4], [2, 1, 2], [2, 3], [3, 2], [4, 1], [5]]
   >>> sums(6, 3)
    [[1, 2, 1, 2], [1, 2, 3], [1, 3, 2], [2, 1, 2, 1], [2, 1, 3], [2, 3, 1], [3, 1, 2],
   [3, 2, 1]]
    0.000
   if n < 0:
        return []
   if n == 0:
                          # The only way to sum to zero using positives
        sums_to_zero = []
        return [sums_to_zero] # Return a list of all the ways to sum to zero
   result = []
   for k in range(1, m + 1):
        result = result + [[k] + rest for rest in sums(n-k, m) if rest == [] or rest[0]
   != k
   return result
```

k is the first number in a list that sums to n, and rest is the rest of that list, so build a list that sums to n.

Call sums to build all of the lists that sum to n-k so that they can be used to construct lists that sum to n by putting a k on the front.

Here is where you ensure that "no two adjacent numbers are the same." Since k will be the first number in the list

you're building, it must not be equal to the first element of rest (which will be the second number in the list you're building).

The recursive case is that each list that sums to n is an integer k (up to m) followed by the elements of a list that sums to n-k and does not start with k.

Here are some key ideas in translating this into code: - If rest is [2, 3], then [1] + rest is [1, 2, 3]. - In the expression [... for rest in sums(...) if ...], rest will be bound to each of the lists within the list returned by the recursive call. For example, if sums(3, 2) was called, then rest would be bound to [1, 2] (and then later [2, 1]).

In the solution, the expresson below creates a list of lists that start with k and are followed by the elements of rest, first checking that rest does not also start with k (which would construct a list starting with two k's).

```
[[k] + rest for rest in sums(n-k, m) if rest == [] or rest[0] != k]
```

For example, in sums (5, 2) with k equal to 2, the recursive call sums (3, 2) would first assign rest to [1, 2], and so [k] + rest would be [2, 1, 2]. Then, it would assign rest to [2, 1] which would be skipped by the if, avoiding [2, 2, 1], which has two adjacent 2's.

You can use [recursion visualizer][RecursionVisualizerSumFun] to step through the call structure of sums(5, 3).

If you get stuck (which many groups do), ask for help!

Optional Question

Tree recursion problems often appear on exams. Here's one:

Q4: A Perfect Question

This question was Fall 2023 Midterm 2 Question 4(a). The original exam version had an extra blank (where total < k * k appears below), but also included some guidance via multiple choice options and hints.

Definition. A perfect square is k*k for some integer k.

Implement fit, which takes positive integers total and n. It returns True or False indicating whether there are n positive perfect squares that sum to total. The perfect squares need not be unique.

```
def fit(total, n):
   """Return whether there are n positive perfect squares that sums to total.
   >>> [fit(4, 1), fit(4, 2), fit(4, 3), fit(4, 4)] # 1*(2*2) for n=1; 4*(1*1) for n=4
    [True, False, False, True]
   >>> [fit(12, n) for n in range(3, 8)] # 3*(2*2), 3*(1*1)+3*3, 4*(1*1)+2*(2*2)
    [True, True, False, True, False]
   >>> [fit(32, 2), fit(32, 3), fit(32, 4), fit(32, 5)] # 2*(4*4), 3*(1*1)+2*2+5*5
    [True, False, False, True]
   def f(total, n, k):
        if total == 0 and n == 0:
            return True
        elif total < k * k:</pre>
            return False
        else:
            return f(total - k * k, n - 1, k) or f(total, n, k + 1)
   return f(total, n, 1)
```